

Last time:

Recovery of individual sparse vecs.  

$$\text{sgn}(z) = \begin{cases} \frac{z}{|z|} & z \neq 0 \\ 0 & z = 0 \end{cases}$$

Thm. 4.26

Given  $A \in \mathbb{C}^{m \times n}$ , the vec.  $x \in \mathbb{C}^n$  with  $\text{supp}(x) = S$  is the unique min. of  $\|z\|_1$  s.t.  $Az = Ax$  if one of the foll. equivalent conditions hold:

- (a)  $|\sum_{j \in S^c} \text{sgn}(x_j) a_j| < \|A_S\|_2$ ,  $\forall v \in \mathcal{N}(A) \setminus \{0\}$   
 (b)  $A_S$  is injective, and  $\exists h \in \mathbb{C}^m$  s.t.  
 $(A^H h)_j = \text{sgn}(x_j)$ ,  $j \in S$   
 $|(A^H h)_j| < 1$ ,  $j \in S^c$ .

Today: - A few remarks about theorem 4.26.  
 - Characterization via tangent cones.

Remarks:

- If  $x \in \mathbb{C}^n$  satisfies (a) of Thm. 4.26, then all vecs.  $x' \in \mathbb{C}^n$  of  $\text{supp}(x') = S' \subset S$  &  $\text{sgn}(x') = \text{sgn}(x)$  are also exactly recovered via BP.
- Thm. 4.26 can be made stable (compressible vecs) and robust (additive noise); see Ex. 4.17 (HW).
- The converse to Thm. 4.26 does not hold in the complex setting, but it does hold in the real-valued case. See Thm. 4.30.

Stable & robust version of the result:

Thm. 4.33

Let  $A = [a_1, \dots, a_n] \in \mathbb{C}^{m \times n}$ , let  $x \in \mathbb{C}^n$  with its largest values supported on  $S$ . Let  $y = Ax + e$  with  $\|e\|_2 \leq \eta$ . For  $\delta, \beta, \gamma, \theta, \tau \geq 0$ ,  $\delta < 1$ , assume

$$\|A_S^H A_S - I\|_2 \leq \delta, \quad \max_{j \in S} \|A_{S^c}^H a_j\|_2 \leq \beta$$

and  $\exists u \in \mathbb{C}^n$  with  $k \in \mathbb{C}^m$  s.t.  
 $\|u_S - \text{sgn}(x_S)\|_2 \leq \gamma$ ,  $\|u_{S^c}\|_\infty \leq \theta$ ,  $\|ku\|_2 \leq \tau\sqrt{\delta}$ .

If  $\beta \leq \theta + \frac{\delta}{1-\delta} < 1$ , then a minimizer  $x^*$  of  $\|z\|_1$  s.t.  $\|Az - y\|_2 \leq \eta$  satisfies:

$$\|x - x^*\|_2 \leq c_1 \|e\|_2 + (c_2 + c_3 \tau) \eta$$

for some const's  $c_1, c_2, c_3 \geq 0$  dep. on  $\delta, \beta, \gamma, \theta, \tau$  as

$$\begin{cases} c_1 = \frac{\delta}{1-\delta} \left(1 + \frac{\beta}{1-\delta}\right); & c_2 = \frac{\delta}{1-\delta} \left(1 + \frac{\beta}{1-\delta}\right) \\ c_3 = \frac{2\sqrt{\delta\theta}}{1-\delta} \left[\frac{\gamma}{1-\delta} \left(1 + \frac{\beta}{1-\delta}\right) + 1\right]. \end{cases}$$

Remark:  $\delta = \beta = \tau = \frac{1}{2}$ ,  $\theta = \frac{1}{4}$ ,  $\tau = 2$   
 $\beta = \frac{1}{2}$ ,  $c_1 = 16$ ,  $c_2 = 10\sqrt{6}$ ,  $c_3 = 32$ .

Next Characterization: Tangent cones to  $\ell_1$  balls:

Given  $x \in \mathbb{R}^n$ , define the convex cone

$$T(x) = \text{cone} \left( \{z - x : z \in \mathbb{R}^n, \|z\|_1 \leq \|x\|_1\} \right)$$

where cone( $T$ ) is the smallest convex cone containing  $T$ . It is defined as:

$$\text{Cone}(T) = \left\{ \sum_{j=1}^m \xi_j x_j : \xi_j \geq 0, x_1, \dots, x_m \in T \right\}$$

Thm. 4.35

For  $A \in \mathbb{R}^{m \times n}$ ,  $x \in \mathbb{R}^n$  is the unique minimizer of  $\|z\|_1$  s.t.  $Az = Ax$  iff  $\mathcal{N}(A) \cap T(x) = \{0\}$ .

Proof: Suppose  $\mathcal{N}(A) \cap T(x) = \{0\}$ .

Let  $x^*$  be an  $\ell_1$  minimizer.

$$\|x^*\|_1 \leq \|x\|_1, \quad Ax^* = Ax.$$

$$\Rightarrow v \in \mathcal{N}(A) \cap T(x) = \{0\}$$

$\Rightarrow x^* = x$ , i.e.,  $x$  is the unique  $\ell_1$  minimizer.

Conversely, suppose  $x$  is the unique  $\ell_1$  minimizer.

$$v \in T(x) \setminus \{0\} \Rightarrow v = \sum \xi_j (x_j - x), \quad \xi_j \geq 0, \|v\|_1 \leq \|x\|_1$$

$$v \neq 0 \Rightarrow \sum \xi_j > 0 \Rightarrow \text{consider } \xi_j' = \frac{\xi_j}{\sum \xi_j}$$

$$\text{If } v \in \mathcal{N}(A), \quad A \left( \sum \xi_j' x_j \right) = Ax \quad (\because \sum \xi_j' = 1)$$

$$\text{while } \left\| \sum \xi_j' x_j \right\|_1 \leq \sum \xi_j' \|x_j\|_1 \leq \|x\|_1$$

By uniqueness of the  $\ell_1$ -min, this  $\Rightarrow \sum \xi_j' x_j = x$ .

$$\Rightarrow v = 0, \text{ a contradiction. Hence, } (T(x) \setminus \{0\}) \cap \mathcal{N}(A) = \emptyset$$

$$\text{or } T(x) \cap \mathcal{N}(A) = \{0\}, \text{ which completes the proof. } \square$$

Can extend to robust recovery:

Thm. 4.36

$A \in \mathbb{R}^{m \times n}$ ,  $x \in \mathbb{R}^n$ ,  $y = Ax + e \in \mathbb{R}^m$

with  $\|e\|_2 \leq \eta$ . If

$$\inf_{v \in T(x)} \|Av\|_2 \geq \tau \text{ for some } \tau > 0, \quad \|v\|_1 = 1$$

then, a minimizer  $x^*$  of  $\|z\|_1$  s.t.  $\|Az - y\|_2 \leq \eta$  satisfies  $\|x - x^*\|_2 \leq \frac{2\eta}{\tau}$ .

Proof: If  $x^* = x$ , nothing to prove. So suppose  $x^* \neq x$ .

$$\text{Then, } \|x^*\|_1 \leq \|x\|_1 \Rightarrow v \in \mathcal{N}(A) \cap T(x) = \{0\}$$

$$\|v\|_2 \geq 1 \Rightarrow \|Av\|_2 \geq \tau \text{ by the assumption}$$

$$\Rightarrow \|A(x^* - x)\|_2 \geq \tau \|x^* - x\|_2$$

$$\text{But } \|A(x^* - x)\|_2 \leq \|A(x^* - y)\|_2 + \|y - Ax\|_2 \leq 2\eta$$

$$\dots \Rightarrow \|x^* - x\|_2 \leq \frac{2\eta}{\tau} \quad \square$$

Hence  $\|z - x\|_2 = \frac{\epsilon}{\sqrt{2}}$ .

Remark: Thm. 4.55, 4.56 extend the above result to the complex case, by defining

$$T(x) = \text{Cone}(\{z - x : z \in \mathbb{C}^n, \|z\|_2 \leq \|x\|_2\}).$$

Low rank matrix recovery

$$X \in \mathbb{C}^{n_1 \times n_2}, \text{rank}(X) \leq r$$

$$y = A(X) \in \mathbb{C}^m$$

$$\downarrow \text{linear map } \mathbb{C}^{n_1 \times n_2} \rightarrow \mathbb{C}^m$$

Equivalent of  $(P_0)$  problem:

$$\min_{z \in \mathbb{C}^{n_1 \times n_2}} \text{rank}(z) \text{ s.t. } A(z) = y$$

This is NP hard (Ex. 2.11).

$$\text{rank}(z) = \|\sigma(z)\|_0 \text{ where } \sigma(z) = \begin{bmatrix} \sigma_1(z) \\ \vdots \\ \sigma_n(z) \end{bmatrix}$$

is a vector containing the singular values of  $z$ .  
Following the  $\ell_1$ -min. philosophy, relax the pb. to:

$$(P_n) : \min_{z \in \mathbb{C}^{n_1 \times n_2}} \|z\|_* \text{ s.t. } A(z) = y$$

$\|\cdot\|_*$  is the nuclear norm:

$$\|z\|_* \cong \sum_{j=1}^n \sigma_j(z), \quad n = \min(n_1, n_2).$$

( $\ell_1$ -norm of the vec. of singular values).

Properties:

- ①  $\|\cdot\|_*$  is a norm (Appendix A.25)
- ②  $(P_n)$  is a convex opt pb.
- ③ Actually equivalent to an SDP.